Theoretical aspects of jet-finding

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Based on work with Cacciari (LPTHE) & Soyez (BNL), and also with Butterworth (ATLAS UCL), Dasgupta (Manchester), Davison (ATLAS UCL), Magnea (Turin), Rojo (LPTHE), Rubin (LPTHE) **Aim:** to provide a reminder/intro to the "basics" of jet-finding, with a couple of bleeding-edge subjects at the end.

General considerations

Some points that deserve to be at the back of your mind as you carry out the challenging technical work to prepare ATLAS to measure jets.

Modern jet definitions

Cone algorithms (and SISCone) Sequential recombination algorithms (k_t , etc.) Learning to compare them

Jets at work

Pileup characterisation and subtraction A new jet-based Higgs search-channel



Jets are what we see. Clearly(?) 2 jets here How many jets do you see? Do you really want to ask yourself this question for 10⁸ events?



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How many jets do you see?

Do you really want to ask yourself this question for 10^8 events?



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- A jet definition is a fully specified set of rules for projecting information from 100's of hadrons, onto a handful of parton-like objects:
 - or project 1000's of calorimeter towers
 - or project dozens of (showered) partons
 - or project a handful of (unshowered) partons

Idea of such a projection: Sterman & Weinberg '77

- Resulting objects (jets) used for many things, e.g. :
 - reconstructing decaying massive particles
 - constraining proton structure
 - as a theoretical tool to attribute structure to an event

E.g. in CKKW matching

e.g. top \rightarrow 3 jets

- ▶ You *lose much information* in projecting event onto jet-like structure:
 - Sometimes information you had no idea how to use
 - Sometimes information you may not trust, or of no relevance

The construction of a jet is unavoidably ambiguous. On at least two fronts:

- 1. which particles get put together into a common jet? Jet algorithm + parameters
- 2. how do you combine their momenta? Recombination scheme Most commonly used: direct 4-vector sums (*E*-scheme)

Taken together, these different elements specify a choice of jetdefinitioncf. Les Houches '07 nomenclature accord

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Physical results (particle discovery, masses, PDFs, coupling) should be independent of your choice of jet definition a bit like renormalisation scale/scheme invariance

Tests independence on modelling of radiation, hadronisation, etc.

- Sometimes there may be a good reason why one jet definition is more optimal than another
 In such cases you should understand why
- But in general there is no single universal 'best' jet definition
 How best to look at an event depends on what you're trying to see

Flexibility in jet-finding is crucial

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"Jetography" like photography



Much fine detail on boarding pass, photograph it from close up, focus=40cm.

keep focus at 40cm

reset focus to 8m

Jets theory, G. Salam (p. 7) Introduction

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Jets should be invariant with respect to certain modifications of the event:

- infrared (IR) emission
 e.g. 1 GeV particles wrt 1 TeV jets
- collinear (C) splitting

This was one of the key ideas introduced by Sterman & Weinberg in "77.

Why?

- ► They happen randomly, quantum-mechanically, all the time in QCD.
- ► IR/C sensitivity ↔ lose real-virtual cancellation in NLO/NNLO QCD calculations → meaningless, divergent results
- ► Hadron-level IR/C modifications are fundamentally non-perturbative
- ▶ Detectors resolve neither full collinear nor full infrared event structure

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The majority of algorithms map each particle into (at most) one jet.

Cone

- top-down
- centred around idea of an 'invariant', directed energy flow

Sequential pairwise recombination $(k_t, \text{ etc.})$

- bottom-up
- successively undoes QCD branching

<u>Others</u>

- "Optimal" Jet Finder (OJF): weight w_{iJ} for each particle *i* to be in jet J
- \blacktriangleright ARCLUS: seq. rec., but with 3 \rightarrow 2 recombination

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Resolve cases of overlapping stable cones

By running a 'split-merge' procedure [Blazey et al. '00 (Run II jet physics)]

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Qu: How do you find the stable cones?

Until recently used iterative methods:

 use each particle as a starting direction for cone; use sum of contents as new starting direction; repeat.

Iterative Cone with Split Merge (IC-SM) e.g. Tevatron cones (JetClu, midpoint) ATLAS cone

- ► Find one stable cone E.g. by iterating from hardest seed particle
- Call it a jet;remove its particles from the event; repeat



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- This is not the same algorithm
- Many physics aspects differ

Iterative Cone with Progressive Removal (IC-PR)

e.g. CMS cone, Pythia Cone, [GetJet], ...

- What are the starting points for iteration?
- Start with hardest particle as seed (IC-PR): collinear unsafe
- Use all particles (IC-SM): extra soft one \rightarrow new solution

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	Last meaningful order			
	ATLAS cone	MidPoint	CMS it. cone	
	[IC-SM]	[IC _{mp} -SM]	[IC-PR]	
Inclusive jets	LO	NLO	NLO	
W/Z + 1 jet	LO	NLO	NLO	
3 jets	none	LO	LO	
W/Z + 2 jets	none	LO	LO	
$m_{ m jet}$ in $2j + X$	none	none	none	

Among consequences of IR unsafety:

NB: 30 - 50M investment in NLO

Some common cone algs (LH '07)

Algorithm	Туре	IRC status	Notes
CDF JetClu	IC _r -SM	IR_{2+1}	
CDF MidPoint cone	IC _{mp} -SM	IR ₃₊₁	
CDF searchcone	IC _{se,mp} -SM	IR_{2+1}	
D0 Run II cone	IC _{mp} -SM	IR ₃₊₁	no seed threshold, but c
ATLAS Cone	IC-SM	IR ₂₊₁	
PxCone	IC _{mp} -SD	IR ₃₊₁	no seed threshold, but c
CMS Iterative Cone	IC-PR	Coll ₃₊₁	
PyCell/CellJet (from Pythia)	FC-PR	Coll ₃₊₁	
GetJet (from ISAJET)	FC-PR	Coll ₃₊₁	

Too many cones, too many problems, needs consolidation.

 $IC-SM \rightarrow SISCone \text{ (finds all stable cones)}$ $IC-PR \rightarrow anti-k_t \text{ [NEW!! More in a few minutes]}$
- 1. Find all distinct ways of enclosing a subset of particles in a $y \phi$ circle
- 2. Check, for each enclosure, if it corresponds to a stable cone

Finding all distinct circular enclosures of a set of points is geometry:



Any enclosure can be moved until a pair of points lies on its edge.

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- 1: Put the set of current particles equal to the set of all particles in the event.
- 2: repeat
- 3: Find *all* stable cones of radius *R* for the current set of particles, e.g. using algorithm 2.
- 4: For each stable cone, create a protojet from the current particles contained in the cone, and add it to the list of protojets.
- 5: Remove all particles that are in stable cones from the list of current particles.
- 6: **until** No new stable cones are found, or one has gone around the loop $N_{\rm pass}$ times.
- 7: Run a Tevatron Run-II type split-merge procedure, algorithm 3, on the full list of protojets, with overlap parameter f and transverse momentum threshold $p_{t,min}$.

- 1: For any group of collinear particles, merge them into a single particle.
- 2: for particle $i = 1 \dots N$ do
- Find all particles *j* within a distance 2*R* of *i*. If there are no such particles, *i* forms a stable cone of its own.
 Otherwise for each *j* identify the two circles for which *i* and *j* lie on the circumference. For each circle, complexity of the complexity
- 4: Otherwise for each j identify the two circles for which i and j lie on the circumference. For each circle, compute the angle of its centre C relative to i, ζ = arctan ^{Δφ_iC}/_{Δν_iC}.
- 5: Sort the circles into increasing angle ζ . 6: Take the first circle in this order, and c
- 6: Take the first circle in this order, and call it the current circle. Calculate the total momentum and checkxor for the cones that it defines. Consider all 4 permutations of edge points being included or excluded. Call these the "current cones".

7: repeat 8: for 9: 10:

for each of the 4 current cones do

If this cone has not yet been found, add it to the list of distinct cones.

If this cone has not yet been labelled as unstable, establish if the in/out status of the edge particles (with respect to the cone momentum axis) is the same as when defining the cone; if it is not, label the cone as unstable.

11: end for 12: Move to

Move to the next circle in order. It differs from the previous one either by a particle entering the circle, or one leaving the circle. Calculate the momentum for the new circle and corresponding new current cones by adding (or removing) the momentum of the particle that has entered (left); the checkxor can be updated by XORing with the label of that particle.

13: until all circles considered.

14: end for

- 15: for each of the cones not labelled as unstable do
- 16: Explicitly check its stability, and if it is stable, add it to the list of stable cones (protojets).
- 17: end for

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Jets theory, G. Salam (p. 17)
Mainstream jet algorithms
Cone
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SISCone part 3: split-merge

1: repeat

Remove all protojets with $p_t < p_{t,\min}$.

Identify the protojet (i) with the highest \tilde{p}_t ($\tilde{p}_{t,jet} = \sum_{i \in jet} |p_{t,i}|$).

Among the remaining protojets identify the one (j) with highest \tilde{p}_t that shares particles (overlaps) with *i*.

- 5: if there is such an overlapping jet then
- 6: Determine the total $\tilde{p}_{t,\text{shared}} = \sum_{k \in i \& j} |p_{t,k}|$ of the particles shared between *i* and *j*.

7: **if**
$$\tilde{p}_{t,\text{shared}} < f \tilde{p}_{t,j}$$
 then

Each particle that is shared between the two protojets is assigned to the one to whose axis it is closest. The protojet momenta are then recalculated.

9: else

Merge the two protojets into a single new protojet (added to the list of protojets, while the two original ones are removed).

11: end if

12: If steps 7–11 produced a protojet that coincides with an existing one, maintain the new protojet as distinct from the existing copy(ies).

13: else

Add *i* to the list of final jets, and remove it from the list of protojets.

- 15: end if
- 16: until no protojets are left.

- Generate event with 2 < N < 10 hard particles, find jets
- Add 1 < N_{soft} < 5 soft particles, find jets again [repeatedly]
- If the jets are different, algorithm is IR unsafe.

Unsafety level	failure rate
2 hard + 1 soft	
3 hard + 1 soft	

Be careful with split-merge too

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Unsafety level	failure rate
2 hard + 1 soft	$\sim 50\%$
3 hard + 1 soft	$\sim 15\%$
SISCone	IR safe !

Be careful with split-merge too



 k_t /Durham algorithm

Majority of QCD branching is soft & collinear, with following divergences:

 $[dk_j]|M_{g\to g_j g_j}^2(k_j)| \simeq \frac{2\alpha_s C_A}{\pi} \frac{dE_j}{\min(E_i, E_j)} \frac{d\theta_{ij}}{\theta_{ij}}, \qquad (E_j \ll E_i, \ \theta_{ij} \ll 1).$

Invert branching process: take pair with strongest divergence between them — they're the most *likely* to belong together.

 \rightarrow k_t/Durham algorithm (e^+e^-)

1. Calculate (or update) distances between all particles *i* and *j*:

$$y_{ij} = rac{2\min(E_i^2, E_j^2)(1 - \cos \theta_{ij})}{Q^2}$$

2. Find smallest of y_{ij}

NB: relative k_t between particles

- ► If > y_{cut}, stop clustering
- Otherwise recombine i and j, and repeat from step 1

Catani, Dokshitzer, Olsson, Turnock & Webber '91

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Exclusive long. inv. k_t algorithm:

- ▶ Drop normalisation to Q^2 (not fixed in pp): $y_{ij} \rightarrow d_{ij}$, $y_{cut} \rightarrow d_{cut}$
- Make it longitudinally boost invariant

 $d_{ij} = \min(p_{ti}^2, p_{tj}^2) \Delta R_{ij}^2, \qquad \Delta R_{ij}^2 = (y_i - y_j)^2 + (\phi_i - \phi_j)^2$

Introduce clustering with "beam jets", distance: d_{iB} = p²_{ti}
 Catani, Dokshitzer, Seymour & Webber '93

Inclusive long. inv. k_t : introduce R, drop d_{cut}

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2. if *ij*, recombine them; *if iB*, *call i a jet* and remove from list of particles

3. repeat from step 1 until no particles left.

S.D. Ellis & Soper, '93; most "cone-like" Jets all separated by at least R on y, ϕ cylinder Often just called " k_t " in pp, $p\bar{p}$ Exclusive long. inv. k_t algorithm:

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k_t is a form of Hierarchical Clustering

Fast Hierarchical Clustering and Other Applications of Dynamic Closest Pairs

David Eppstein UC Irvine

We develop data structures for dynamic closest pair problems with arbitrary distance functions, that do not necessarily come from any geometric structure on the objects. Based on a technique previously used by the author for Euclidean closest pair, we show how to insert and delete objects from an nobject set, maintaining the doest pair, in $O(\log^2 n)$ time per update and O(n) pace. With quadratic space, we can instead use a quadtree-like structure to achieve an optimal time bound, O(n) per update. We apply these data structures to hierarchical clustering, greedy matching, and TSP heuristics, and discuss other potential applications in machine learning. Gröbner bases, and local improvement algorithms for partition and placement problems. Experiments show our new methods to be faster in practice than previously used heuristics.

Categories and Subject Descriptors: F.2.2 [Analysis of Algorithms]: Nonnumeric Algorithms

General Terms: Closest Pair, Agglomerative Clustering

Additional Key Words and Phrases: TSP, matching, conga line data structure, quadtree, nearest neighbor heuristic

1. INTRODUCTION

Hierarchical clustering has long been a mainstay of statistical analysis, and clustering based methods have attracted attention in other fields: computational biology (reconstruction of evolutionary trees; tree-based multiple sequence alignment), scientific simulation (n-body problems), theoretical computer science (network design and nearest neighbor searching) and of course the web (hierarchical indices such as Yahoo). Many clustering methods have been devised and used in these applications, but less effort has gone into algorithmic speedups of these methods.

In this paper we identify and demonstrate speedups for a key subroutine used in several clustering algorithms, that of maintaining closest pairs in a dynamic set of objects. We also describe several other applications or potential applications of the Idea behind k_t alg. is to be found over and over in many areas of (computer) science.

Sequential recombination of cacti

kt alg.: Find smallest of

 $d_{ij} = \min(k_{ti}^2, k_{tj}^2) \Delta R_{ij}^2 / R^2, \quad d_{iB} = k_{ti}^2$

If d_{ij} recombine; if d_{iB} , *i* is a jet Example clustering with k_t algorithm, R = 0.7

























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If d_{ij} recombine; if d_{iB} , *i* is a jet Example clustering with k_t algorithm, R = 0.7

Sequential recombination of cacti



k_t alg.: Find smallest of $d_{ij} = \min(k_{ti}^2, k_{tj}^2) \Delta R_{ij}^2 / R^2, \quad d_{iB} = k_{ti}^2$

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 $d_{ij} = \min(k_{ti}^{2p}, k_{tj}^{2p}) \Delta R_{ij}^2 / R^2$ $d_{iB} = k_{ti}^{2p}$

	Alg. name	Comment	time
p = 1	k _t	Hierarchical in rel. k_t	
	CDOSTW '91-93; ES '93		NIn N exp.
p = 0	Cambridge/Aachen	Hierarchical in angle	
	Dok, Leder, Moretti, Webber '97	Scan multiple <i>R</i> at once	N In N
	Wengler, Wobisch '98	$\leftrightarrow QCD \text{ angular orderin}$	
p = -1	anti- k_t Cacciari, GPS, Soyez '08	Hierarchy meaningless.	
	$\sim {\sf reverse}{-}k_t$ Delsart, Loch et al.	Behaves like IC-PR	$N^{3/2}$
SC-SM	SISCone	Replacement for IC-SM	
	GPS Soyez '07 + Tevatron run II '00	notably "MidPoint" cones	$N^2 \ln N \exp$.

One could invent/try others (e.g. OJF, etc.). Our [Paris+BNL] philosophy: 4 algs is enough of a basis to develop first physics understanding.

We already have far more than can be shown here

non-COMMERCIAL BREAK

One place to stop for your jet-finding needs:

FastJet

http://www.lpthe.jussieu.fr/~salam/fastjet Cacciari, GPS & Soyez '05-08

- Fast, native, computational-geometry methods for k_t, Cam/Aachen, anti-k_t
 Cacciari & GPS '05-06
- Plugins for SISCone (plus some other, deprecated, legacy cones)
- Documented user interface for adding extra algorithms of your own
- ► Tools for jet areas, pileup characterisation & subtraction

Physics quality measures

E.g. width of narrowest window (around mass-peak) containing $\sim 25\%$ of events that pass basic selection cuts, in hadronic $t\bar{t}$ and Z' events.



Les Houches: Cacciari, Rojo, GPS & Soyez '08

Reach of jet algorithms



Herwig 6.510 + FastJet 2.1

Jet contours - visualised









<u>LEP</u>

- $M_{BSM} \sim 1 \text{ TeV}$
- ▶ *M_{EW}* ~ 100 GeV
- $p_{t,\text{pileup}} \sim 25 100 \text{ GeV/unit rap.}$
- ▶ $p_{t,UE} \sim 2.5 5 \text{ GeV/unit rap.}$
- $p_{t,hadr.} \sim 0.5 \text{ GeV/unit rap.}$

Multitude of scales forces us to go beyond "1 parton = 1 jet" equivalence

 $\sim \alpha_{s} \times M_{BSM}$

 $\sim M_{EW}$

 $\sim lpha_{\sf s} imes M_{EW}$

<u>Tevatron</u>

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Multitude of scales forces us to go beyond "1 parton = 1 jet" equivalence

 $\sim M_{FW}$

 $\sim \alpha_{s} \times M_{RSM}$







Add dense coverage of infinitely soft *"ghosts"* See how many end up in jet to measure jet area

 \sim 10000 particles Clustering takes \sim 0.2s



Jet areas



Jet areas



Median p_t /area across the set of jets in an event is a good estimator of pileup+UE in *that event*

Area-based subtraction

Basic Procedure:

- Use p_t/A from majority of jets (pileup jets) to get level, ρ, of pileup and UE in event
- Subtract pileup from hard jets:

$$p_t \rightarrow p_{t,sub} = p_t - A\rho$$

Cacciari & GPS '07

Illustration:

- semi-leptonic $t\bar{t}$ production at LHC
- high-lumi pileup (\sim 20 ev/bunch-X)

Same simple procedure works for a range of algorithms



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Jets theory, G. Salam (p. 34) └──Taking jets further └──A new Higgs search-channel?

Searching for high- p_t HW/WZ?

High- p_t light Higgs decays to $b\bar{b}$ inside a single jet. Can this be seen? Butterworth, Davison, Rubin & GPS '08



Then on the Higgs-candidate: *filter* away UE/pileup by reducing $R \rightarrow R_{filt}$, take *three hardest subjets* (keep LO gluon radⁿ) + require *b*-tags on two hardest.

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Jets theory, G. Salam (p. 35) └──Taking jets further └──A new Higgs search-channel?

Compare with "standard" algorithms

Check mass spectra in HZ channel, $H \rightarrow b\bar{b}$, $Z \rightarrow \ell^+ \ell^-$



Cambridge/Aachen (C/A) with mass-drop and filtering (MD/F) works best

Jets theory, G. Salam (p. 36) Taking jets further A new Higgs search-channel?

combine HZ and HW, $p_t > 200 \text{ GeV}$

Combined HZ + HW, pt > 200 GeV



- ► Take $Z \to \ell^+ \ell^-$, $Z \to \nu \bar{\nu}$, $W \to \ell \nu$ $\ell = e, \mu$
- ▶ p_{tV}, p_{tH} > 200 GeV
- ▶ $|\eta_V|, |\eta_H| < 2.5$
- Assume real/fake *b*-tag rates of 0.7/0.01.
- Some extra cuts in HW channels to reject tt.
- Assume $m_H = 115$ GeV.

At 5.9σ for 30 fb⁻¹ this looks like a competitive channel for light Higgs discovery. **Deserves serious exp. study!**

- ► Know the algorithms and make sure they're fully described somewhere
- ▶ Be sure they're infrared & collinear safe (there's plenty of choice)
- LHC's a very different environment from what we've had before
- One should explore a range of jet definitions for any given analysis
- Both to optimize the analysis
- And to verify robustness of conclusions wrt jet-def
- Jets are not just partons
- Tell you lots about UE, pileup \rightarrow model it better, subtract it better
- Have substructure, which can help reveal underlying physics
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EXTRA MATERIAL



Extra soft particle adds new seed \rightarrow changes final jet configuration.

This is **IR unsafe**.

Divergences of real and virtual contributions do not cancel at $\mathcal{O}\left(\alpha_{\rm s}^4\right)$

Kilgore & Giele '97

Solution: add extra seeds at midpoints of all pairs, triplets, . . . of stable cones. Seymour '97 (?)



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Midpoint cone alg. misses some stable cones; extra soft particle \rightarrow extra starting point \rightarrow extra stable cone found **MIDPOINT IS INFRARED UNSAFE**

Or collinear unsafe with seed threshold



 $\label{eq:misses} \begin{array}{l} \mbox{Midpoint cone alg. misses some stable cones; extra soft} \\ \mbox{particle} \rightarrow \mbox{extra starting point} \rightarrow \mbox{extra stable cone found} \\ \mbox{MIDPOINT IS INFRARED UNSAFE} \end{array}$

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Midpoint cone alg. misses some stable cones; extra soft particle \rightarrow extra starting point \rightarrow extra stable cone found **MIDPOINT IS INFRARED UNSAFE**

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Compare midpoint and SISCone

Result depends on observable:

- inclusive jet spectrum is the least sensitive (affected at NNLO)
- larger differences (5 10%) at hadron level

seedless reduces UE effect





Jets theory, G. Salam (p. 42) Extras IC-SM unsafety

Look at jet masses in multijet events. NB: Jet masses reconstruct boosted W/Z/H/top in BSM searches



 $\begin{array}{l} \mbox{Select 3-jet events} \\ p_{t1,2,3} > \{120,60,20\} \mbox{ GeV}, \end{array}$

Calculate LO jet-mass spectrum for jet 2, compare midpoint with SISCone.

▶ 10% differences by default

► 40% differences with extra cut △R_{2,3} < 1.4 e.g. for jets from common decay chain

In complex events, IR safety matters

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Jets theory, G. Salam (p. 43)
Extras
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Showering puts in many extra seeds: missing stable cones (in midpoint) should be less important?

Look at 3rd jet mass distribution (no ΔR_{23} cut):



Missing stable cones \rightarrow 50% effects even after showering

There are N(N-1)/2 distances d_{ij} — surely we have to calculate them all in order to find smallest?

kt distance measure is partly geometrical:

- Consider smallest $d_{ij} = \min(k_{ti}^2, k_{tj}^2)R_{ij}^2$
- Suppose $k_{ti} < k_{tj}$
- ▶ Then: $R_{ij} \leq R_{i\ell}$ for any $\ell \neq j$. [If $\exists \ \ell \text{ s.t. } R_{i\ell} < R_{ij}$ then $d_{i\ell} < d_{ij}$]

In words: if i, j form smallest d_{ij} then j is geometrical nearest neighbour (GNN) of i.

 k_t distance need only be calculated between GNNs

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Jets theory, G. Salam (p. 45)

Extras

Speed
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Finding Geom Nearest Neighbours



Given a set of vertices on plane (1...10) a *Voronoi diagram* partitions plane into cells containing all points closest to each vertex Dirichlet '1850. Voronoi '1908

A vertex's nearest other vertex is always in an adjacent cell.

E.g. GNN of point 7 will be found among 1,4,2,8,3 (it turns out to be 3)

Construction of Voronoi diagram for *N* points: *N* In *N* time Fortune '88 Update of 1 point in Voronoi diagram: In *N* time Devillers '99 [+ related work by other authors]

Convenient C++ package available: **CGAL** http://www.cgal.org Assemble with other comp. science methods: **FastJet** Cacciari & GPS, hep-ph/0512210 http://www.lpthe.jussieu.fr/~salam/fastjet/

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Status in 2005





Status in 2007



Single package, FastJet, to access all developments, natively (k_t, Cam/Aachen) or as plugins (SISCone): Cacciari, GPS & Soyez '05-08 http://www.lpthe.jussieu.fr/~salam/fastjet/

	Jet $\langle \delta p_t \rangle$ given by product of dependence on				
	scale	colour factor	R	\sqrt{s}	
pert. radiation	$\sim rac{lpha_{\sf s}({\sf p}_t)}{\pi}{\sf p}_t$	Ci	$\ln R+\mathcal{O}\left(1\right)$	-	
hadronisation	Λ_h	Ci	$-1/R + \mathcal{O}\left(R ight)$	_	
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To get best experimental resolutions, minimise contributions from all 3 components.

> Here: sum of squared means Better still: calculate flucts

NB: this is rough picture details of p_t scaling wron

But can still be used to understand general principles.

Dasgupta, Magnea & GPS '07

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Optimal $R \vee p_t$, proc., collider



This kind of information is the start of what might go into **"auto-focus for jetography"**



Optimal $R \vee p_t$, proc., collider



This kind of information is the start of what might go into "auto-focus for jetography" Correlate MC & measured UE (LHC)



Check subtraction



How can high- p_t HV do better than low p_t search?

- Acceptance improves at high-p_t
 - H and V both more likely to be central
 - H and V, since boosted, have central decay products
 - Decay products have high p_t, so you always see them
- Backgrounds fall faster than signal
 - for $\nabla b\bar{b}$, not gigantic effect (factor 2–3), but there
 - ► for $t\bar{t}$, top mass \rightarrow b's with $p_t = 60 70$ GeV, just like light Higgs at rest; much harder to fake high- p_t H
- ▶ At high- p_t capture most of gluon radiation from $H \rightarrow b\bar{b} + X$, without too much UE; much harder at low p_t . → Better mass resolution.

